

Online Learnability and Complexity Measures

Project Presentation - CMPUT 654, ML Theory

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What if instead of “batch” data, we have a stream of data?

Online Learning: Random Averages, Combinatorial Parameters,
and Learnability

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Batch (Offline)

PAC Learning



Finite VC Dim,
Rademacher,...

- Existence of an efficient algorithm finding a Probably Approximately Correct hypothesis.

w.p. $1 - \delta$:

$$err_D(\hat{f}) - err_D(f^*) \leq \epsilon$$

- Generic algorithm:
Empirical Risk Minimization (ERM)
- Oracle inequalities

Online

Sublinear
Regret



Finite LDim, Seq.
Rademacher,....

- Regret:

$$\frac{1}{T} \left(\sum_{t=1}^T err_t(\hat{f}_t) - err_t(f^*) \right) \in o(1)$$

- Generic algorithm:
Standard Optimal Algorithm (SOA)

- N. Littlestone, Learning quickly when irrelevant attributes abound: A new linear-threshold algorithm. 1988

Online Learning

Hypothesis Class \mathcal{F} Feature Set \mathcal{X}

A learner interacts with an adversary over T rounds.

In each round $t = 1, 2, \dots, T$:

Learner picks a distribution $q_t \in \mathcal{Q}$ over functions $f: \mathcal{X} \rightarrow \mathcal{Y}$

Adversary picks a feature-label pair x_t

Learner draws a sample $f_t \sim q_t$ and suffers loss $f_t(x_t)$

Repetitive Two-Player
Game

What's its value?

Regret:

$$\mathbb{E} \inf_{f \in \mathcal{F}} \sum_{t=1}^T [f_t(x_t) - f(x_t)]$$

Online Learning

Hypothesis Class \mathcal{F}

Feature Set \mathcal{X}

Label Set $\mathcal{Y} \subseteq \mathbb{R}$

A learner interacts with an adversary over T rounds.

In each round $t = 1, 2, \dots, T$:

Learner picks a distribution $q_t \in \mathcal{Q}$ over functions $f: \mathcal{X} \rightarrow \mathcal{Y}$

Adversary picks a feature-label pair (x_t, y_t)

Learner draws a sample $f_t \sim q_t$ and suffers loss $err_t(f_t) = \ell(f_t(x_t) - y_t)$

Repetitive Two-Player
Game

What's its value?

Regret:

$$\mathbb{E} \sum_{t=1}^T \left[err_t(f_t) - \inf_{f^* \in \mathcal{F}} err_t(f^*) \right]$$

Realizable: $\exists f^* : f^*(x_t) = y_t \quad \forall t$

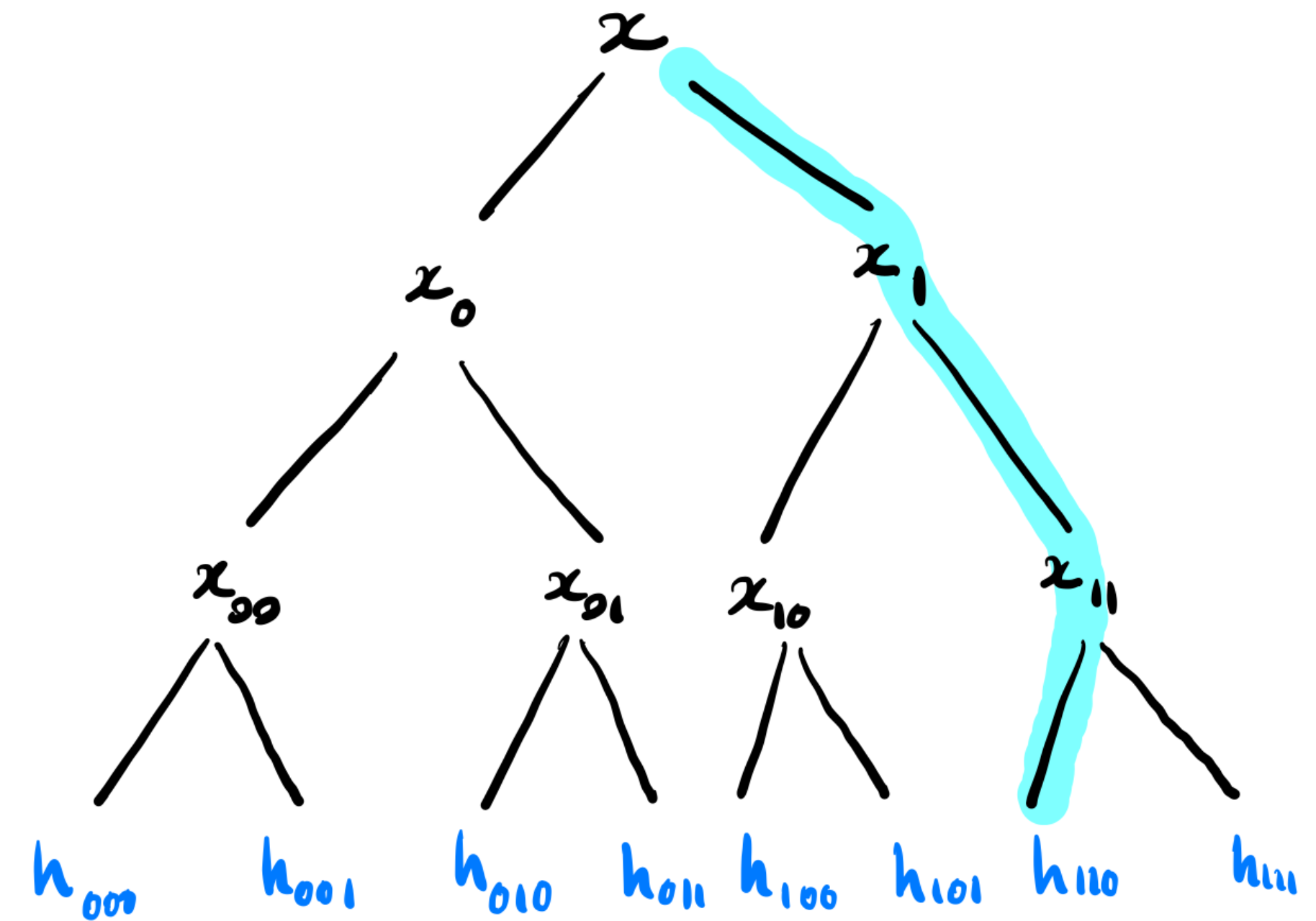
Agnostic: o.w.

Littlestone's Dimension

- Analogous to VC dimension (Binary)

Shattered Tree: \forall path $\epsilon, \exists f \in \mathcal{F}$

Ldim := Max depth of a shattered labeled tree.



Definition 6. An \mathcal{X} -valued tree \mathbf{x} of depth d is *shattered* by a function class $\mathcal{F} \subseteq \{\pm 1\}^{\mathcal{X}}$ if for all $\epsilon \in \{\pm 1\}^d$, there exists $f \in \mathcal{F}$ such that $f(\mathbf{x}_t(\epsilon)) = \epsilon_t$ for all $t \in [d]$. The *Littlestone dimension* $\text{Ldim}(\mathcal{F}, \mathcal{X})$ is the largest d such that \mathcal{F} shatters an \mathcal{X} -valued tree of depth d .

- Continuous version?! fat-shattering

- N. Littlestone, Learning quickly when irrelevant attributes abound: A new linear-threshold algorithm. 1988
 - S. Ben David et al, Agnostic online learning, 2009

Online Learnability

- Value of the Game:

$$\mathcal{V}_T(\mathcal{F}, \mathcal{X}) = \inf_{q_1 \in \mathcal{Q}} \sup_{x_1 \in \mathcal{X}} \mathbb{E}_{f_1 \sim q_1} \cdots \inf_{q_T \in \mathcal{Q}} \sup_{x_T \in \mathcal{X}} \mathbb{E}_{f_T \sim q_T} \left[\sum_{t=1}^T f_t(x_t) - \inf_{f \in \mathcal{F}} \sum_{t=1}^T f(x_t) \right]$$

Prokhorov's theorem

$$= \sup_{p_1} \mathbb{E}_{x_1 \sim p_1} \cdots \sup_{p_T} \mathbb{E}_{x_T \sim p_T} \left[\sum_{t=1}^T \inf_{f_t \in \mathcal{F}} \mathbb{E}_{x_t \sim p_t} [f_t(x_t)] - \inf_{f \in \mathcal{F}} \sum_{t=1}^T f(x_t) \right]$$

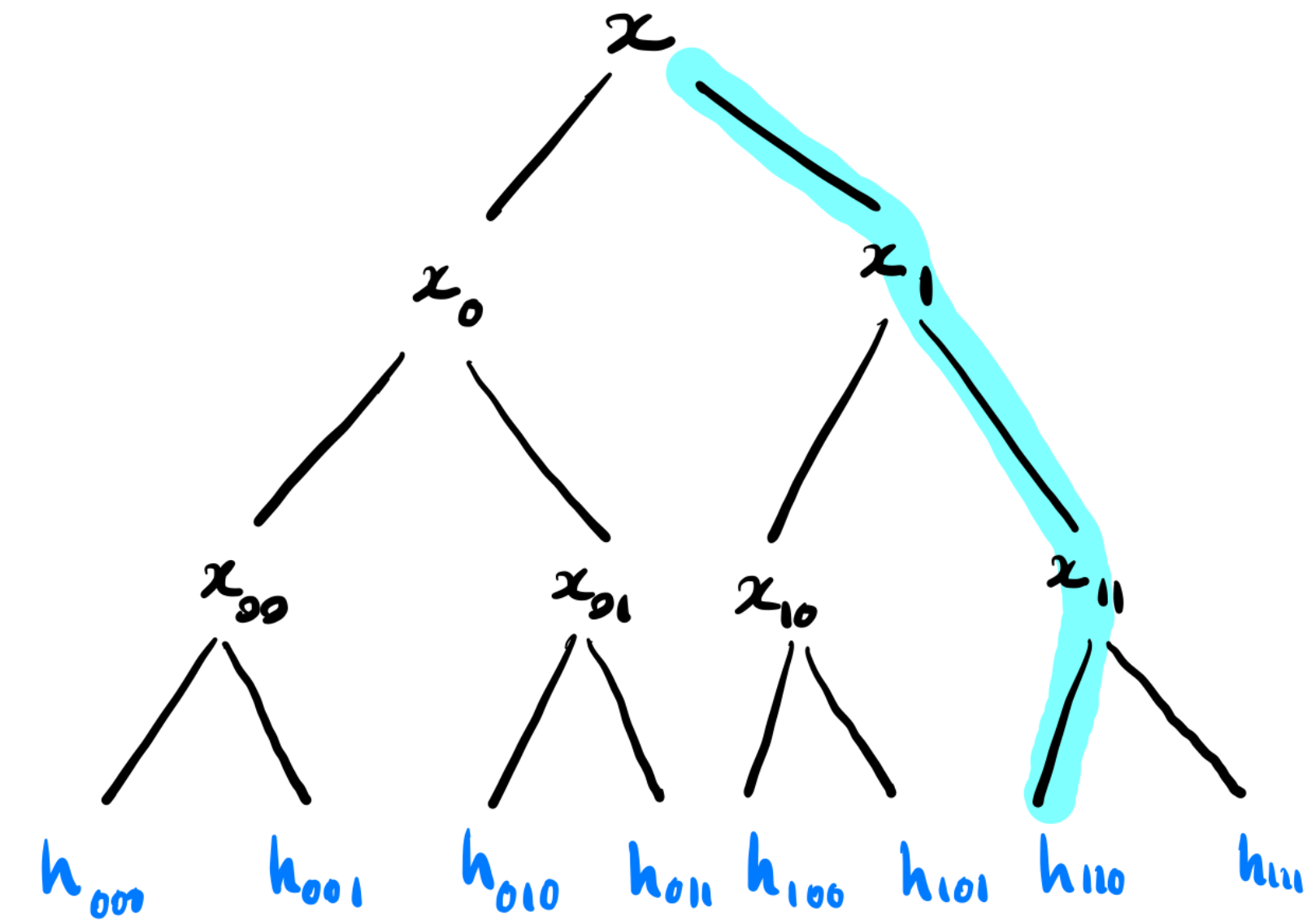
- \mathcal{F} is “online learnable” w.r.t \mathcal{X} if:

$$\limsup_{T \rightarrow \infty} \frac{\mathcal{V}_T(\mathcal{F}, \mathcal{X})}{T} = 0 .$$

Sequential Rademacher Complexity

- Like offline, but ϵ and samples are no longer independent:

$$\mathfrak{R}_T(\mathcal{F}) = \sup_{\mathbf{x}} \mathbb{E}_{\epsilon} \left[\sup_{f \in \mathcal{F}} \sum_{t=1}^T \epsilon_t f(\mathbf{x}_t(\epsilon)) \right]$$



- Minimax value is bounded with Sequential Rademacher complexity:

$$\mathcal{V}_T(\mathcal{F}) \leq 2\mathfrak{R}_T(\mathcal{F})$$

- A bit more tricky symmetrization...

Standard Optimal Algorithm (SOA)

- ERM of online learning
- Compute the little stone dimension of a sub-hypothesis class

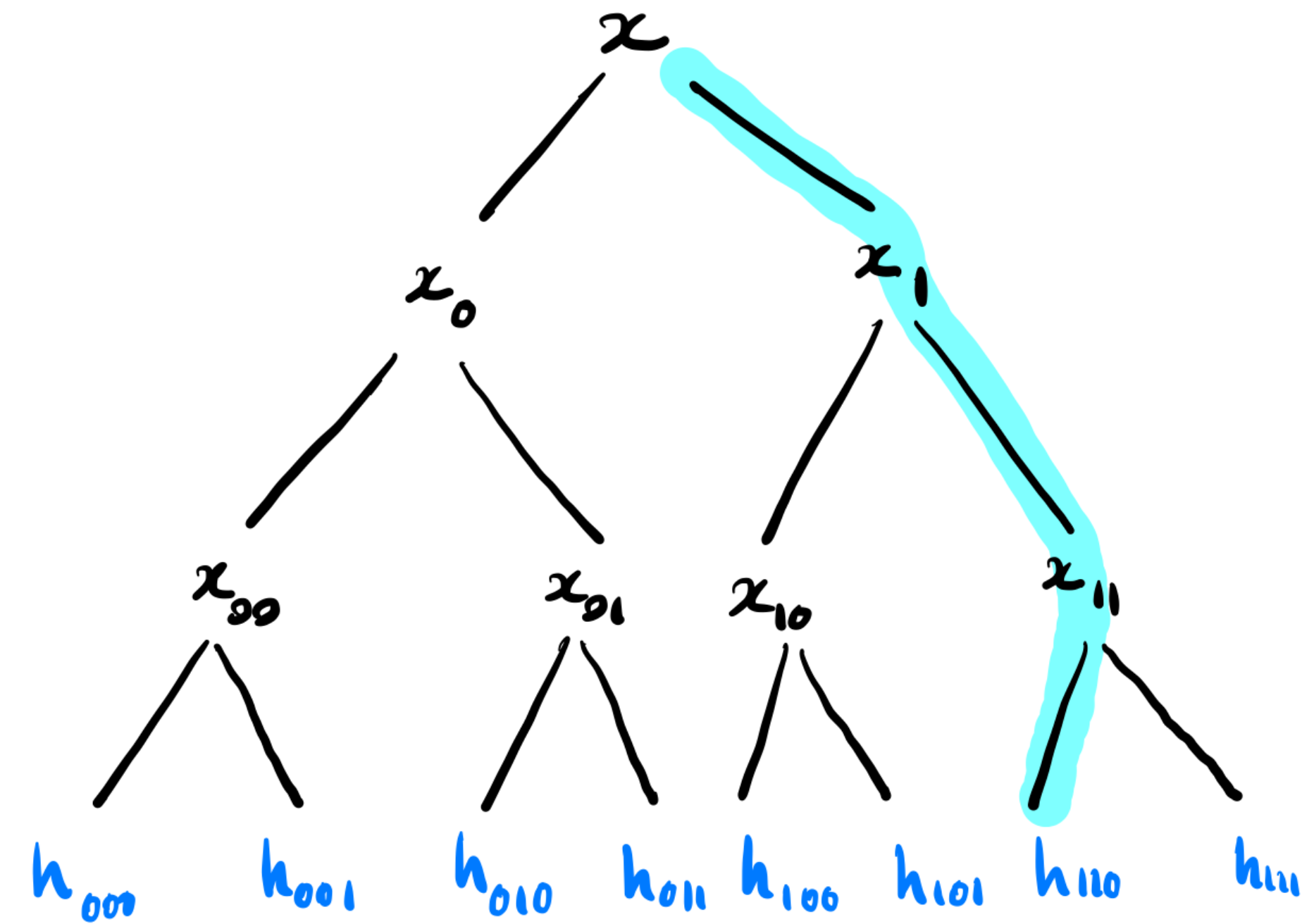
- Realizable setting:

$$Reg \leq Ldim$$

- Agnostic setting:

$$Reg \leq \tilde{O}(\sqrt{Ldim \cdot T})$$

- Making T^{Ldim} many experts s.t. \exists optimal experts \rightarrow Expert Advise problem



What if I use ERM in Online Learning?

- Not too sublinear!
- Still need exponentially number of queries to ERM

Online Learning and Solving Infinite Games with an ERM Oracle

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Unapproximability of Littlestone's Dimension

- Hard to approximate!

Improved Inapproximability of VC Dimension and Littlestone's Dimension via (Unbalanced) Biclique

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Thank you